

Generalising Kim's lemma to abstract independence relations

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Workshop on Model Theory and Applications (Chania)

Outline

- 1 Background
- 2 Generalising Kim's lemma
- 3 Consequences

What is an independence relation?

Fix a complete theory T and monster $\mathbb{M} \models T$.

For us, an **independence relation** is an $\text{Aut}(\mathbb{M})$ -invariant ternary relation on small subsets of $\mathbb{M}$. We denote it by $\downarrow$, and write $A \downarrow_C B$ if $(A, B, C) \in \downarrow$.

For the rest of this talk, every independence relation $\downarrow$ will satisfy the following two properties:

- *Full existence*: for any sets A, B and C , there is $A' \equiv_C A$ such that $A' \downarrow_C B$.
- *Monotonicity*: if $A \downarrow_C B$ and $A' \subseteq A$, $B' \subseteq B$, then $A' \downarrow_C B'$.

A look at the classification map

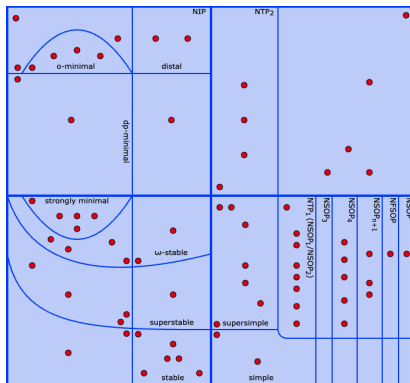


Figure: forkinganddividing.com (Conant)

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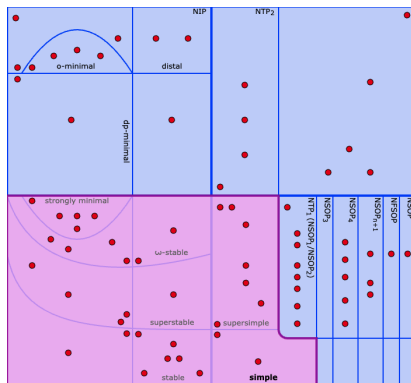


Figure: forkinganddividing.com (Conant)

Simple theories

Fact (Kim's Lemma)

A theory T is simple iff, whenever $\varphi(x, b)$ divides over C , it divides along every Morley sequence over C starting at b .

Kim and Pillay characterised simplicity in terms of the behaviour of non-forking:

Definition

$A \downarrow_C^f B :=$ there is no formula in $\text{tp}(A/BC)$ that forks over C .

If T is simple:

- $\downarrow^f$ is symmetric.
- $\downarrow^f$ satisfies local character.
- $\downarrow^f$ satisfies the independence theorem over models.

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Adler's framework

Definition

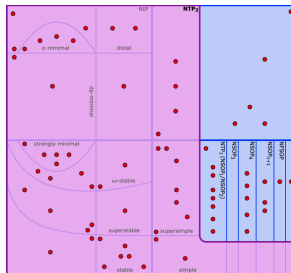
We say $\perp$ is an **Adler independence relation** if it satisfies monotonicity, right base monotonicity, left transitivity, left normality, finite character, local character, and right extension.

Each of these properties is “algebraic”.

Fact (Adler)

Every Adler independence relation is symmetric.

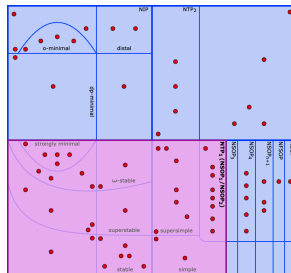
Beyond simplicity



NTP_2 (Chernikov-Kaplan):

- A theory T is NTP_2 iff, whenever $\varphi(x, b)$ divides over C , it divides along every *strict* Morley sequence over C starting at b .
- If T is NTP_2 , forking and dividing coincide over models.

Beyond simplicity



NSOP₁ (Kaplan, Shelah, Ramsey):

- A theory T is NSOP₁ iff, whenever $\varphi(x, b)$ Kim-divides over M , it divides along every M -invariant Morley sequence starting at b .
- If T is NSOP₁, $\perp^K$ satisfies symmetry, *chain* local character, and the independence theorem over models.

[illegible]

- In many NSOP₄ theories, there is a stationary independence relation satisfying the *relative Kim's lemma*.
- If *Conant*-independence is symmetric, then T is NSOP₄.

Main question

How can we understand the recent developments in neostability theory in terms of abstract independence relations?

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Generalised Morley sequences

Fix an independence relation $\perp$. We look at “special” sequences:

Definition

A sequence $(a_i)_{i < \omega}$ is $\perp$ -**Morley over** C if it is C -indiscernible and $a_i \perp_C a_{<i}$ for all $i < \omega$.

Examples

- $(a_i)_{i < \omega}$ is $\perp^f$ -Morley over C iff it is Morley over C .
- $(a_i)_{i < \omega}$ is $\perp^i$ -Morley over M iff it is M -invariant Morley.
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The Key Idea

For an independence relation $\perp$:

- $\varphi(x, b) \perp$ -**Kim-divides over** M if it divides along *some* $\perp$ -Morley sequence over M .
- $\varphi(x, b) \perp$ -**Conant-divides over** M if it divides along *every* $\perp$ -Morley sequence over M .
- $\varphi(x, b) \perp$ -**Kim-forks** (resp., $\perp$ -**Conant-forks**) **over** M if $\varphi(x, b) \vdash \bigvee_{i < n} \psi_i(x, c_i)$ and each $\psi_i(x, c_i) \perp$ -Kim-divides (resp., $\perp$ -Conant-divides) over M .

Freeness

We say $\varphi(x, b)$ is $\perp$ -**free over** C if there is some $a \models \varphi(x, b)$ such that $a \perp_C b$.

Examples

- $\varphi(x, b)$ is $\perp^u$ -free over C iff there is $c \in C$ such that $\models \varphi(c, b)$.
- $\varphi(x, b)$ is $\perp^f$ -free over C iff it does not fork over C .
- $\varphi(x, b)$ is $\perp^a$ -free over C iff it does not *strongly fork* over C .

$A \perp_C^u B$ iff $\text{tp}(A/BC)$ is a coheir over C

$A \perp_C^a B$ iff $\text{acl}(AC) \cap \text{acl}(BC) \subseteq \text{acl}(C)$

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Some Adlerian operations

We also need some operations on an independence relation $\perp$:

- Define $\perp^*$ by “forcing right extension”: $A \perp_C^* B$ iff, for every $B' \supseteq B$, there is $A' \equiv_{BC} A$ such that $A' \perp_C B'$.
- Define $\perp^{\text{sfc}}$ by “forcing strong finite character”: $A \perp_C^{\text{sfc}} B$ iff every formula in $\text{tp}(A/BC)$ is $\perp$ -free over C .

Relativising Kim- and Conant-independence

- Define $A \perp_M^+ B$ iff, for every $\perp$ -Morley sequence I over M starting at B , there is $A' \equiv_{MB} A$ such that I is MA' -indiscernible.

• $\perp$ -Kim-independence = $((\perp^+)^{\text{sfc}})^*$.

- We say $\varphi(x, b)$ **weakly $\perp$ -Kim-divides over M** if it is *not* $\perp^+$ -free over M .

- We say $\varphi(x, b)$ **$\perp$ -Kim-forks over M** if it is *not* $(\perp^+)^*$ -free over M .

- Define $A \perp_M^\times B$ iff there exists some MA -indiscernible $\perp$ -Morley sequence over M starting at B .

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Relativising Kim- and Conant-independence

Theorem

- $\varphi(x, b) \perp\!\!\!\perp$ -Conant-divides over M iff it divides along every $\perp\!\!\!\perp$ -Morley sequence over M .
- $\varphi(x, b) \perp\!\!\!\perp$ -Kim-forks (resp., $\perp\!\!\!\perp$ -Conant-forks) over M iff $\varphi(x, b) \vdash \bigvee_{i < n} \psi_i(x, c_i)$ and each $\psi_i(x, c_i) \perp\!\!\!\perp$ -Kim-divides (resp., $\perp\!\!\!\perp$ -Conant-divides) over M .
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Back to Kim's lemma

We will say $\perp^1$ **has GUWP w.r.t.** $\perp^2$ if, whenever $\varphi(x, b)$ $\perp^2$ -Kim-divides over M , it also $\perp^1$ -Conant-divides over M .

Theorem

*Suppose $\perp^1$ and $\perp^2$ satisfy full existence and monotonicity.
TFAE:*

- ① $\perp^1$ witnesses $\perp^2$.
- ② If $\varphi(x, b)$ weakly $\perp^2$ -Kim-divides over M , then it also $\perp^1$ -Conant-divides over M .
- ③ $((\perp^1)^\times)^{\text{sfc}} \implies ((\perp^2)^+)^{\text{sfc}}$.

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Back to Kim's lemma

Examples

- If T is simple, then $\perp^f$ witnesses $\perp^a$.
- If T is NTP_2 , then $\perp^{\text{ist}}$ witnesses $\perp^a$.
- If T is NSOP_1 , then $\perp^i$ witnesses $\perp^i$.
- If T is NBTP , then $\perp^{K^{\text{st}}}$ witnesses $\perp^i$.
- If T is NCTP , then $\perp^{\text{bi}}$ witnesses $\perp^i$.

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The Old

Theorem

Suppose $\perp$ satisfies full existence and monotonicity. Assume $\perp$ is “compatible” with $\perp$ -Kim-independence, and $\perp$ -Kim-independence satisfies left transitivity. TFAE:

- ① $\perp$ -Kim-independence is symmetric.
- ② $\perp$ -Kim-independence witnesses $\perp$.

In this case, $\perp$ -Kim-independence also satisfies chain local character.

Examples

- ① This recovers symmetry and local character for $\perp^f$ in simple theories.
- ② Assuming symmetry, we recover “witnessing” and chain local character for Kim-independence in NSOP₁ theories.

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Suppose $\downarrow' \Rightarrow \downarrow$ satisfy full existence, monotonicity and right extension, and additionally $\downarrow'$ satisfies left extension and $\downarrow$ “weak transitivity”. Assume $\downarrow \Rightarrow \downarrow$ -Kim-independence. If $\downarrow$ -Kim-independence satisfies left transitivity and symmetry, then it also satisfies the $\downarrow'$ -independence theorem over models.

Examples

- 1 Independence theorem over models in simple theories.
- 2 Assuming symmetry, “weak independence theorem” in NSOP_1 theories.
- 3 $\downarrow$ -independence theorem for $\downarrow$ -Kim-independence whenever $\downarrow$ is a *free amalgamation relation*.

The Old

Theorem

Suppose $\perp' \Rightarrow \perp$ satisfy full existence, monotonicity and right extension, and additionally $\perp'$ satisfies left extension and $\perp$ “weak transitivity”. Assume $\perp \Rightarrow \perp$ -Kim-independence. If $\perp$ -Kim-independence satisfies left transitivity and symmetry, then it also satisfies the $\perp'$ -independence theorem over models.

Examples

- 1 Independence theorem over models in simple theories.
- 2 Assuming symmetry, “weak independence theorem” in NSOP_1 theories.
- 3 $\perp$ -independence theorem for $\perp$ -Kim-independence whenever $\perp$ is a *free amalgamation relation*.

The New

Definition (repeated)

We say $\perp$ is an **Adler independence relation** if it satisfies monotonicity, right base monotonicity, left transitivity, left normality, finite character, local character, and right extension.

Theorem

For any Adler independence relation $\perp$:

- 1 $\perp = \perp^\times$. In particular, $\varphi(x, b)$ is $\perp$ -free over M iff it does not $\perp$ -Conant-fork over M .
- 2 $\perp$ satisfies strong finite character iff $\perp = \perp$ -Conant-independence.

The New

Theorem

Suppose $\perp$ satisfies full existence, monotonicity, and quasi-strong finite character¹. If $\perp^i$ witnesses $\perp$, then T is either NSOP₁ or BTP.

As a consequence, this shows that most strictly NSOP₄ theories have BTP.

Conjecture

If T is NSOP₄, then it is either NSOP₁ or BTP.

¹That is, for types $p(x), q(y) \in S(M)$, " $p(x) \cup q(y) \cup x \perp_M y$ " is type-definable over M .

The Old and New

Theorem

Let $\perp^1 \Rightarrow \perp^2$ be independence relations satisfying full existence and monotonicity such that $\perp^1$ also satisfies left and right extension. If $\perp^2$ -Conant-independence satisfies the $\perp^2$ -independence theorem over models, then $\perp^1$ witnesses $\perp^2$.

This recovers, without using trees, the fact that T is NSOP₁ iff Conant-independence satisfies the independence theorem over models (Chernikov, Kaplan, Ramsey).

The Old and New

Theorem (Mutchnik)

Let $\perp$ be an independence relation satisfying full existence, monotonicity, stationarity, and the generalised freedom axiom. If $\perp$ -Kim-independence is symmetric, then $\perp^i$ witnesses $\perp$.

Proof.

Using the properties of $\perp$, it can be shown that $\perp$ -Kim-independence satisfies left transitivity. Thus, by our Theorem, as $\perp$ -Kim-independence is symmetric, it follows that $\perp$ -Kim-independence witnesses $\perp$. But $\perp^i \implies \perp$ -Kim-independence. Thus, $\perp^i$ witnesses $\perp$. □

The End

Thank you!